

Decomposition of the Complete r -Graph into Complete r -Partite r -Graphs*

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Abstract. For $n \geq r \geq 1$, let $f_r(n)$ denote the minimum number q , such that it is possible to partition all edges of the complete r -graph on n vertices into q complete r -partite r -graphs. Graham and Pollak showed that $f_2(n) = n - 1$. Here we observe that $f_3(n) = n - 2$ and show that for every fixed $r \geq 2$, there are positive constants $c_1(r)$ and $c_2(r)$ such that $c_1(r) \leq f_r(n) \cdot n^{-\lfloor r/2 \rfloor} \leq c_2(r)$ for all $n \geq r$. This solves a problem of Aharoni and Linial. The proof uses some simple ideas of linear algebra.

1. Introduction

For $n \geq r \geq 1$, let $f_r(n)$ denote the minimum number q , such that it is possible to partition all edges of the complete r -uniform hypergraph on n vertices into q pairwise edge-disjoint complete r -partite r -uniform hypergraphs.

Obviously, $f_1(n) = 1$. Graham and Pollak ([3, 4], see also [2, 5]) proved that $f_2(n) = n - 1$ for all $n \geq 2$. Simple proofs for this result were found by Tverberg [7] and Peck [6].

Aharoni and Linial [1] raised the natural problem of determining or estimating $f_r(n)$ for $r > 2$. In particular they asked if $f_r(n)$ is a nonlinear function of n , for some fixed $r > 2$.

In this note we answer this question in the affirmative by proving the following theorem, that determines the asymptotic behavior of $f_r(n)$ for every fixed r as n tends to infinity.

Theorem 1.1. *For every fixed $r \geq 1$, there are two positive constants $c_1 = c_1(r)$ and $c_2 = c_2(r)$ such that*

$$c_1 \cdot n^{\lfloor r/2 \rfloor} \leq f_r(n) \leq c_2 \cdot n^{\lfloor r/2 \rfloor}$$

for all $n \geq r$.

The lower bound is proved using some simple ideas of linear algebra. The method is similar to the one used by Tverberg [7] and by Graham and Pollak [3, 4], for determining $f_2(n)$. The upper bound is established by a recursive construction.

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It is worth noting that our construction supplies the exact value of $f_3(n) = n - 2$ for all $n \geq 3$.

2. The Lower Bound

We start with the following easy observation.

Lemma 2.1. *For every $n \geq r \geq 2$*

$$f_r(n) \geq f_{r-1}(n-1).$$

Proof. Suppose all edges of the complete r -uniform hypergraph on a set $N = \{1, 2, \dots, n\}$ of n vertices are partitioned into $q = f_r(n)$ r -partite r -graphs (= r uniform hypergraphs) H^1, H^2, \dots, H^q . Let E_i denote the set of edges of H^i and put $\bar{E}_i = \{e - \{n\} : e \in E_i, n \in e\}$. Clearly each nonempty \bar{E}_i is the set of edges of a complete $(r-1)$ -partite $(r-1)$ -graph. Moreover, the set of all nonempty \bar{E}_i 's forms a decomposition of all edges of the complete $r-1$ -uniform hypergraph on the $n-1$ vertices $N - \{n\}$. Hence $f_{r-1}(n-1) \leq q = f_r(n)$, as needed. \square

In view of Lemma 2.1, the lower bound in Theorem 1.1 for odd values of r follows from the lower bound for even values of r , which we prove next.

Lemma 2.2. *For all $n \geq 2k \geq 2$*

$$f_{2k}(n) \geq 2 \cdot \frac{\binom{n}{k} - \binom{n}{k-1} - \binom{n}{k-3} - \dots - \binom{n}{k+1-2 \cdot \lceil k/2 \rceil}}{\binom{2k}{k}}.$$

Proof. Let $\underline{K} = \{K \subset N : |K| = k\}$ be the set of all $\binom{n}{k}$ k -subsets of $N = \{1, 2, \dots, n\}$ and associate each $K \in \underline{K}$ with a variable x_K . Let H be a complete $2k$ -partite $2k$ -graph, whose (pairwise disjoint) vertex classes V_1, V_2, \dots, V_{2k} are subsets of N . By definition, the edges of H are all $2k$ -subsets $A \subset N$, such that $|A \cap V_i| = 1$ for $1 \leq i \leq 2k$. We define, for each such H , a quadratic form $Q(H)$ in the variables $\{x_K : K \in \underline{K}\}$ as follows.

$Q(H) = \sum \{L_A(H) \cdot L_B(H) : A, B \subset \{1, 2, \dots, 2k\}, |A| = |B| = k, A \cap B = \emptyset, 1 \in A\}$, where, for $C \subset \{1, 2, \dots, 2k\}, |C| = k$,

$$L_C(H) = \sum \{x_K : K \in \underline{K}, |K \cap V_c| = 1 \text{ for all } c \in C\}.$$

Thus, $Q(H)$ is a sum of $\frac{1}{2} \binom{2k}{k}$ products of the form $L_A(H) \cdot L_B(H)$, in which each factor is a linear combination of the x_K 's.

Put $q = f_{2k}(n)$, and suppose the edges of the complete r -graph on N are partitioned into q r -partite r -graphs H^1, H^2, \dots, H^q . One can easily check that

$$\sum_{i=1}^q Q(H^i) = \sum \{x_K \cdot x_L : K, L \in \underline{K}, K \cap L = \emptyset\}. \quad (2.1)$$

Indeed, if $K, L \in \underline{K}$ and $K \cap L = \emptyset$ then the product $x_K \cdot x_L$ appears only in $Q(H^i)$,

where H^i is the unique H^j containing $K \cup L$ as an edge, and if $K \cap L \neq \emptyset$, then $x_k \cdot x_L$ appears in no $Q(H^i)$.

We next claim that

$$\begin{cases} \sum \{x_K \cdot x_L : K, L \in \underline{K}, K \cap L = \emptyset\} \\ = \frac{1}{2} \sum_{i=0}^k (-1)^i \sum_{A \subset N, |A|=i} \left(\sum_{K \in \underline{K}, A \subset K} x_K \right)^2. \end{cases} \quad (2.2)$$

Indeed, if $K, L \in \underline{K}$ and $|K \cap L| = j$, ($0 \leq j < k$), then the coefficient of $x_K \cdot x_L$ in the right hand side of (2.2) is $\sum_{i=0}^j (-1)^i \binom{j}{i}$, which is 1 if $j = 0$ and 0 if $j > 0$. For

$K = L$, the coefficients of x_K^2 in the right hand side of (2.2) is $\frac{1}{2} \sum_{i=0}^k (-1)^i \binom{k}{i} = 0$.

Thus (2.2) holds.

Substituting (2.2) and the definition of the $Q(H^i)$'s into (2.1) we conclude that

$$\begin{cases} \sum_{i=1}^q \sum \{L_A(H^i) \cdot L_B(H^i) : A, B \subset \{1, \dots, 2k\}, |A| = |B| = k, A \cap B = \emptyset, 1 \in A\} \\ = \frac{1}{2} \sum_{i=0}^k (-1)^i \sum_{A \subset N, |A|=i} \left(\sum_{K \in \underline{K}, A \subset K} x_K \right)^2 \end{cases} \quad (2.3)$$

Let V be the linear subspace of the real $\binom{n}{k}$ -dimensional space of the x_K 's determined by the following set of

$$\frac{1}{2} \binom{2k}{k} \cdot q + \binom{n}{k-1} + \binom{n}{k-3} + \dots + \binom{n}{k+1-2\lceil k/2 \rceil}$$

linear equations.

$$\begin{cases} L_A(H^i) = 0 \text{ for all } 1 \leq i \leq q \text{ and } A \subset \{1, 2, \dots, 2k\}, |A| = k, 1 \in A. \\ \sum_{K \in \underline{K}, A \subset K} x_K = 0 \text{ for all } A \subset N, |A| \in \{k-1, k-3, \dots, k+1-2\lceil k/2 \rceil\} \end{cases} \quad (2.4)$$

We claim that V is the zero subspace. Indeed, suppose $\{\bar{x}_K : K \in \underline{K}\} \in V$. Then \bar{x}_K satisfies (2.4), and in view of (2.3) we conclude that

$$0 = \frac{1}{2} \cdot (-1)^k \left\{ \sum_{K \in \underline{K}} \bar{x}_K^2 + \sum_{A \subset N, |A|=k-2} \left(\sum_{K \in \underline{K}, A \subset K} \bar{x}_K \right)^2 + \dots \right\},$$

and hence $\bar{x}_K = 0$ for all $K \in \underline{K}$.

Therefore, the number of linear equations in the system (2.4) is at least $\binom{n}{k}$ and the assertion of Lemma 2.2 follows. \square

Combining Lemmas 2.1 and 2.2 we obtain

Corollary 2.3. For every fixed $r \geq 1$, $f_r(n) \geq c_r \cdot n^{\lceil r/2 \rceil} \cdot (1 + o(1))$ as $n \rightarrow \infty$, where

$$c_r = \frac{2\lceil r/2 \rceil!}{(2\lceil r/2 \rceil)!}.$$

Remarks.

- 1) Lemma 2.2 with $k = 1$ reduces to Graham-Pollak's result; $f_2(n) \geq n - 1$, (which is, of course, sharp).
- 2) Lemma 2.1 with $r = 3$ asserts $f_3(n) \geq f_2(n - 1) = n - 2$. As shown in the next section this result is also sharp.
- 3) A trivial lower bound for $f_r(n)$ is $f_r(n) \geq \binom{n}{r} / \left(\frac{n}{r}\right)^r$, since the number of edges of any complete r -partite r -graph on n vertices is not greater than $(n/r)^r$. This trivial bound is much weaker than the one proved above for all $r = o(n)$, but is better for, e.g., $r = \lfloor n/2 \rfloor$.

3. The Upper Bound

In this section we prove the upper bound for $f_r(n)$ given in Theorem 1.1, using some simple recursive constructions. We first determine $f_3(n)$ for all $n \geq 3$.

Lemma 3.1. *For all $n \geq 3$*

$$f_3(n) = n - 2.$$

Proof. By Lemma 2.1 and Graham-Pollak's result

$$f_3(n) \geq f_2(n - 1) = n - 2.$$

We prove that $f_3(n) \leq n - 2$ by induction on n . For $n = 2, 3$ the result is trivial. Assuming the result for all $n', n' < n$, we prove it for n , ($n > 3$). Put $N = \{1, 2, \dots, n\}$ and $N_i = \{2i - 1, 2i\}$ for $1 \leq i \leq \lfloor n/2 \rfloor$. For odd n define also $N_{\lfloor n/2 \rfloor} = \{n\}$. We claim that

$$f_3(n) \leq \lfloor n/2 \rfloor + f_3(\lceil n/2 \rceil). \quad (3.1)$$

Indeed, put $q = f_3(\lceil n/2 \rceil)$ and let H^1, \dots, H^q be a decomposition of the complete 3-graph on $\lceil n/2 \rceil$ vertices $\{1, 2, \dots, \lceil n/2 \rceil\}$ into q complete 3-partite 3-graphs. For $1 \leq i \leq q$, let V_1^i, V_2^i and V_3^i denote the vertex-classes of H^i . Let \bar{H}^i be the 3-partite 3-graph whose vertex classes are $\cup \{N_j; j \in V_1^i\}$, $\cup \{N_j; j \in V_2^i\}$ and $\cup \{N_j; j \in V_3^i\}$. For $1 \leq j \leq \lfloor n/2 \rfloor$, let \bar{H}^{q+j} be the 3-partite 3-graph whose vertex classes are $\{2i - 1\}$, $\{2i\}$ and $N - \{2i - 1, 2i\}$. One can easily check that the hypergraphs $\{\bar{H}^i\}_{i=1}^{q+\lfloor n/2 \rfloor}$ form a decomposition of all edges of the complete 3-graph on N into 3-partite 3-graphs. This establishes (3.1). Hence, by the induction hypothesis,

$$f_3(n) \leq \lfloor n/2 \rfloor + f_3(\lceil n/2 \rceil) \leq \lfloor n/2 \rfloor + \lceil n/2 \rceil - 2 = n - 2. \quad \square$$

Let N_1 and N_2 be two disjoint sets of vertices, and let H_i be an r_i -graph on N_i , ($1 \leq i \leq 2$). We denote by $H_1 + H_2$ the $(r_1 + r_2)$ -graph on $N_1 \cup N_2$ whose edges are all edges $e_1 \cup e_2$, where e_i is an edge of H_i ($i = 1, 2$). One can easily check that if H_i is a complete r_i -partite r_i -graph then $H_1 + H_2$ is a complete $(r_1 + r_2)$ -partite $(r_1 + r_2)$ -graph. For notational convenience let us agree that $f_0(n) = 1$ for all n .

Lemma 3.2. *Suppose $n \geq r \geq 4$, then*

$$f_r(n) \leq \sum_{i=0}^{r-1} f_i(\lfloor n/2 \rfloor) \cdot f_{r-i}(\lceil n/2 \rceil).$$

Proof. Put $N = \{1, 2, \dots, n\}$, $N_1 = \{1, 2, \dots, \lfloor n/2 \rfloor\}$, $N_2 = \{\lfloor n/2 \rfloor + 1, \dots, n\}$. For $0 \leq i \leq r$ let \underline{H}^i be a family of $f_i(\lfloor n/2 \rfloor)$ complete i -partite i -graphs that decompose the complete i -graph on N_1 . (\underline{H}^0 consists of one graph whose only edge is the empty edge.) Similarly, let \underline{G}^j be a family of $f_j(\lceil n/2 \rceil)$ complete j -partite j -graphs that decompose the complete j -graph on N_2 ($0 \leq j \leq r$). Define a family \underline{F} of $\sum_{i=0}^r f_i(\lfloor n/2 \rfloor) \cdot f_{r-i}(\lceil n/2 \rceil)$ complete r -partite r -graphs on N by

$$\underline{F} = \bigcup_{i=0}^r \{H^i + G^{r-i} : H^i \in \underline{H}^i, G^{r-i} \in \underline{G}^{r-i}\}.$$

One can easily check that the members of \underline{F} form a decomposition of the complete r -graph on N . This completes the proof. \square

We can now prove the upper bound for $f_r(n)$ given in Theorem 1.1 by double induction on r and n . Since $f_r(n)$ is a monotone increasing function of n , it is enough to prove it when n is a power of 2, which we assume, for convenience. By Lemma 3.1 (and trivial constructions for $r \leq 2$) $f_r(n) \leq c_r \cdot n^{\lfloor r/2 \rfloor}$ for $r = 0, 1, 2, 3$ and every n , where $c_0 = c_1 = c_2 = c_3 = 1$. Clearly, if $n < r$ then $f_r(n) \leq c_r \cdot n^{\lfloor r/2 \rfloor}$ for every positive c_r . Assuming that

$$f_{r'}(n') \leq c_{r'} \cdot n'^{\lfloor r'/2 \rfloor} \quad (3.2)$$

for all $r' < r$ and all $n' = 2^j$, and for $r' = r$ and $2^j = n' < n = 2^i$, we prove that if c_r is properly chosen then (3.2) holds also for (r, n) . Indeed, by Lemma 3.2 and the induction hypothesis

$$\begin{aligned} f_r(n) &\leq \sum_{i=0}^r f_i(n/2) f_{r-i}(n/2) \leq \sum_{i=0}^r c_i \cdot c_{r-i} (n/2)^{\lfloor i/2 \rfloor + \lfloor (r-i)/2 \rfloor} \\ &\leq \sum_{i=0}^r c_i \cdot c_{r-i} \cdot (n/2)^{\lfloor r/2 \rfloor} = \frac{1}{2^{\lfloor r/2 \rfloor}} \left(2c_r + \sum_{i=1}^{r-1} c_i \cdot c_{r-i} \right) n^{\lfloor r/2 \rfloor}. \end{aligned}$$

Hence, if we define the c_r 's by

$$\begin{aligned} c_0 = c_1 = c_2 = c_3 = 1 \\ \text{and } c_r \cdot (2^{\lfloor r/2 \rfloor} - 2) = \sum_{i=1}^{r-1} c_i \cdot c_{r-i} \text{ for } r \geq 4, \end{aligned} \quad (3.3)$$

then $f_r(n) \leq c_r \cdot n^{\lfloor r/2 \rfloor}$ for every r and every $n = 2^i$. This implies the validity of the upper bound for $f_r(n)$ given in Theorem 1.1.

Remarks.

1) One can easily check that the constants $\{c_r\}_{r=0}^{\infty}$ defined by (3.3) satisfy

$$c_r \leq \frac{8^{r-1}}{\lfloor r/2 \rfloor!}.$$

By a somewhat more careful analysis we can show that the construction described above implies that for every fixed $k \geq 1$ $f_{2k}(n) \leq \frac{1}{k!} \cdot n^k (1 + o(1))$, as $n \rightarrow \infty$. This should be compared to the lower bound

$$f_{2k}(n) \geq \frac{2k!}{(2k)!} \cdot n^k(1 + o(1))$$

given in Corollary 2.3.

2) It would be interesting to determine $f_r(n)$ precisely for $r > 3$, or to improve our estimates. In particular, Lemma 2.2 and Lemma 3.2 for $r = 4$ imply that

$$\frac{1}{6}(n^2 - 3n) \leq f_4(n) \leq \frac{1}{2}(n^2 - 5n + 6)$$

for all n . It would be interesting to decide which of these two bounds is closer to the truth.

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